Consensus Mechanism Design based on Structured DAGs

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HKUST

Blockchain



WHAT IS GREAT

✓ Secure: Nakamoto consensus

✓ Decentralized: PoW

WHAT NEEDS IMPROVEMENT

TPS

- □ Latency: waiting + confirmation
- **C** Concentration of mining power
- **Transactions** with little fees











Directed Acyclic Graphs (DAG)

5

Block Size in bytes		
id_{prev}	32	
id_{ms}	32	
id _{tip}	32	
nonce	4	
peer	65	
message	~ 500	

Peer chain

Previous block by the same peer

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Peer chain

Previous block by the same peer

Nakamoto chain

Milestone block



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Connectivity

Tip block by another peer



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Header **Overhead** = Header + Message

*			
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4		id_{ms}	32
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			65
		L	00
		message	~ 500
1 T×	2 Tx's	3 Tx's	4 Tx's
17%	9%	6%	5%

Proof of Work

 $B = (id_{prev}, id_{ms}, id_{tip}, nonce, message)$

 $H(B) = 0 \cdots 0 * * * * * * *$

leading x bits are all 0

 $H(B) = 0 \cdots 0 * * * *$

leading y bits are all 0

Regular Block

regular block reward + Tx fee

Milestone Block

additional bonus*

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Three pointers have to be specified before knowing the type!





Peer chain



Peer chain

Nakamoto chain



Peer chain

Nakamoto chain





Peer chain

Nakamoto chain





Peer chain





Level Set



Level Set

 $\mathcal{S}(\mathsf{B}_m,\mathsf{B}_m')=\mathcal{C}(\mathsf{B}_m)\setminus\mathcal{C}(\mathsf{B}_m')$





Peers obtain different parts of the big block from different peers continuously over time.

Transaction Assignment





Miners

Transaction Assignment





Miners

Transaction Assignment





Miners



$H(\text{miner's state}, \text{Tx}) < c \times \text{miner's hashing power}$



Transaction Assignment waste Tx (\$1) (\$<u>100</u> Tx (\$100) T× (\$10) Tx (\$2) (\$100 Tx (\$5) Miners Mempool

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verifiable



 q_i consensus

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 q_i consensus

Protocol Receiving a block

Suppose Alice has a local DAG

$$\mathcal{G}_a = \mathcal{C}(\mathbf{B}_m)$$

- 1. Download if height is too small
- 2. Solidify + topological sort
- 3. Add blocks one by one

Protocol RECEIVING A BLOCK

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Bob's local DAG



Bob's local DAG

Alice's local DAG





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syntactically valid

- What if not pointing to the miner's previous block
- What if not pointing to the most recent milestone
- □ What if not pointing to a recent regular block by others



Post-order DFS



Post-order DFS

- ✓ chronological order of blocks on a peer chain
- ✓ approximate chronological order of all blocks

Order all transactions in a level set Tx_1, Tx_2, Tx_3, \ldots



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 $\mathcal{L}_{k+1} = \mathcal{L}_{k+1} \cup (\mathtt{Tx}_{k+1}),$ $\mathcal{U}_{k+1} = \mathcal{U}_k \cup \{ \text{outputs from } \mathtt{Tx}_{k+1} \} \setminus \{ \text{inputs from } \mathtt{Tx}_{k+1} \}$



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- ✓ No peer id/sig to save message space
- ✓ No coinbase to reduce the number of UTXOs

Туре	Status	Transaction	Reward	Bonus
regular +	on peer chain	Valid	r + Tx fee	0
regular +	on peer chain	Invalid	r	0
regular +	forked	Valid	0	0
regular +	forked	Invalid	0	0
milestone	longest MS chain	Valid	r + Tx fee	$\% \times r \times level set$
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Mot to fork peer chain

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If the second se

Model Analysis

 $H(miner's state, Tx) < c \times miner's hashing power$

 $\mathbb{P}(\text{Tx is not workable for any miner}) \approx e^{-c}$

Model Analysis

Waste

$$\theta(c) = \frac{(1 - e^{-\mu \overline{t}})\mu c\overline{t}}{1 + (1 - e^{-\mu \overline{t}})\mu c\overline{t}}$$

Queueing latency

$$\frac{1}{c} \frac{1}{\rho\mu} \ln \left(\frac{1}{1 - \frac{\rho}{1 - \theta(c)}} \right)$$

Infection delay

$$\frac{2+2\ln(n)}{\mu} + \frac{1}{np\mu}$$



Performance

System parameters

Partition factor c	Block generation speed	MS interval	Avg. # of blocks per level set
0.01	1200 blocks/second	10 seconds	12000



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Modeling assumptions

Tx arrival rate	Percentage of malicious hashing power
1000	30%



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Tx arrival rate	Percentage of malicious hashing power
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Queueing latency	Infection delay	Security latency	Wasted capacity		
188 seconds	23 seconds	810 seconds	1.7%		



Modeling assumptions

Performance

Summary

consensus

reward

algorithm

- ☑ TPS
- In Latency: waiting + confirmation
- Concentration of mining power
- **I** Transactions with little fees

Summary

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hm economic model

- □ How to issue new coins: inflation, limited supply, …?
- How to incentivize people to provide
 storage and bandwidth service?

Summary

consensus

reward

algorithm

- TPS
- Latency: waiting + confirmation $\mathbf{\underline{\vee}}$
- Concentration of mining power
- Transactions with little fees

economic model

optimization

- □ How to issue new coins: inflation, limited supply, ...?
- □ How to incentivize people to provide storage and bandwidth service?
- How to adaptively adjust the number of blocks created per unit of time?

Consensus Mechanism Design based on Structured Directed Acyclic Graphs

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Abstract

We introduce a structure for the directed acyclic graph (DAG) and a mechanism design based on that structure so that peers can reach consensus at large scale based on proof of work (PoW). We also design a mempool transaction assignment method based on the DAG structure to render negligible the probability that a transaction being processed by more than one miners. The result is a significant scale-up of the capacity without sacrificing security and decentralization.

Key words: consensus; directed acyclic graph; proof of work; transaction assignment.



arXiv:<u>1901.02755</u>

GitHub

coming soon